# Theory of Modelling and Simulation

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# Outline

#### $oldsymbol{X}$ Motivation

**★** Hierarchies of System Specification

#### **✗** System morphisms

- **X** System morphisms at the same level of specification
- **X** System morphisms between different levels of specification
- **★** Analysis and verification
- $\pmb{\mathsf{X}}$  Category Theory: the meta-theory





## Motivation

- Modelling: description of a (dynamic) system
- Simulation: generation of possible behaviours of a system given a model
- Analysis: reasoning about a system's behaviour



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# Motivation

- Formalisms for
  - modelling: DEVS, Statecharts, Petri Nets, ODEs, PDEs, etc.
  - analysis: temporal logics, other logics.
  - simulation: state trajectories
- Theory:
  - Foundation
  - Common framework to explain/describe formalisms involved
  - Tools for analysis: general properties of systems and/or formalisms
  - Basis for (software) tools



MSD

## Progress

 $\checkmark$  Motivation

- **★** Hierarchies of System Specification
- **✗** System morphisms
  - **X** System morphisms at the same level of specification
  - $\pmb{\times}$  System morphisms at different levels of specification
- **★** Analysis and verification
- **★** Category Theory: the meta-theory





- Model = system specification
- A model can be given at different levels of *abstraction* or *specification*
- Modelling process (by refinement:) from abstract to concrete



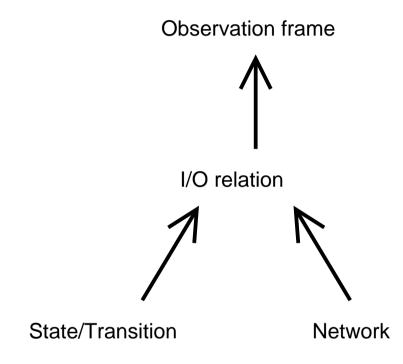
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- Dynamic systems: time-varying behaviour
- *Time-base*: ordered set (with a few properties.)
- *Signals* (or *trajectories*): functions from the time-base to some set
- *Segment*: function from a time-interval to some set



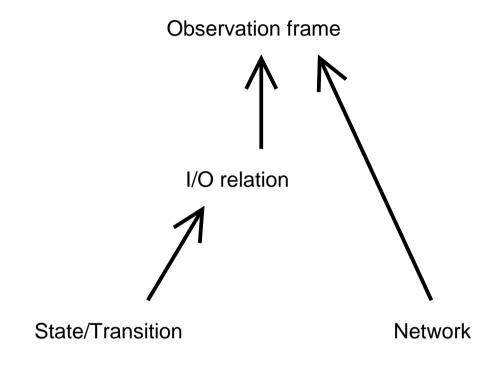
- A simple hierarchy:
  - Observation frame: inputs and outputs
  - I/O relation
  - State/Transition
  - Network





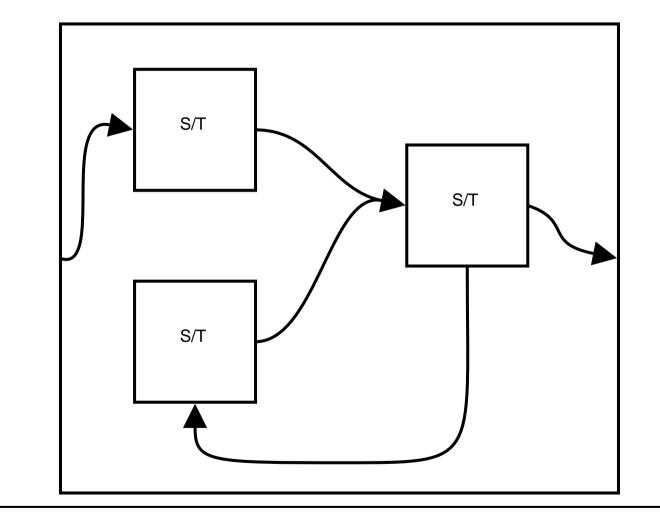




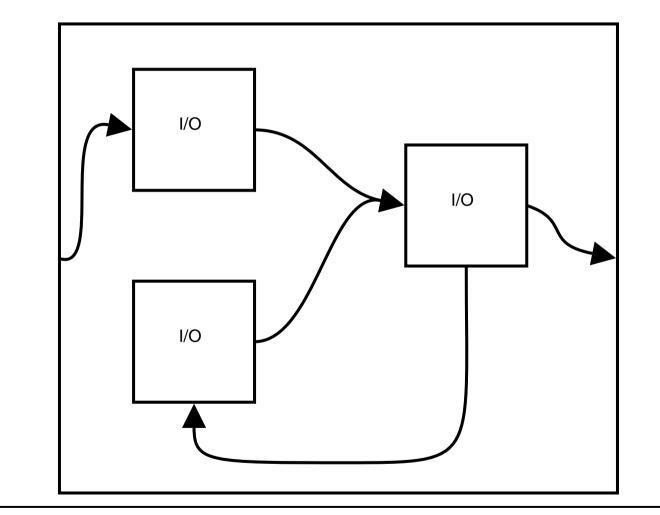




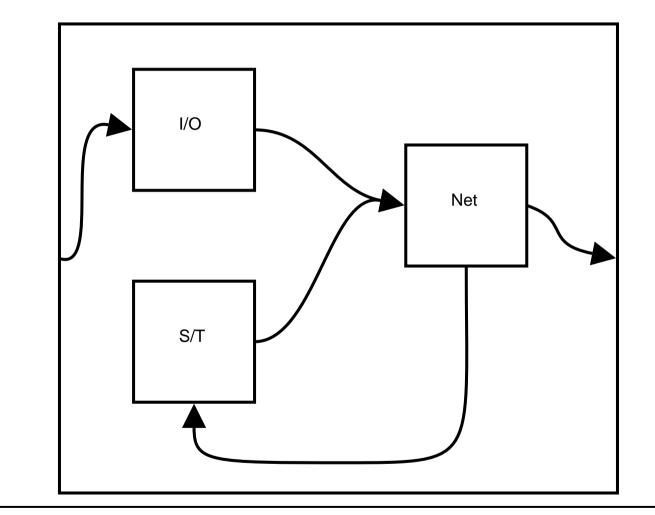














- A general hierarchy of system specification is not a total order but a partial order
- Different kinds of abstraction relationships:
  - Structural
  - Behavioural
- Fitting multiple formalisms in the hierarchy
- Mapping formalisms into the hierarchy
  - Formalisms semantics
  - Multi-formalism modelling



## Progress

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- ✔ Hierarchies of System Specification
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- $\pmb{\times}$  Analysis and verification
- **★** Category Theory: the meta-theory



## System morphisms

- Some questions:
  - Can I plug-in this component in that network?
  - Can I put system A in place of system B and obtain the same behaviour?
  - I want my system to have this I/O relation. Does this system satisfy it?
  - I know system A has this property. Does system B have that property as well?
- Relate different systems



- At the Observation frame level:
  - Same time-base
  - Same input and output sets
  - Same interface



- At the Observation frame level:
  - Equivalent time-base
  - Equivalent input and output sets
  - Equivalent interface



- At the Observation frame level:
  - Compatible time-base
  - Compatible input and output sets
  - Compatible interface
- Existence of a map between
  - The time-bases (speed)
  - The input and output sets
  - The interfaces



- At the I/O relation level
  - Same I/O relation
  - Containment
  - Bijection
  - Bijection + transformation



- At the S/T level
  - What does it mean for a state/transition system to behave in the same way as another?
  - The concepts of simulation and bisimulation



- A labelled transition system (LTS) is a tuple  $(S, L, \rightarrow)$  where
  - S is a set of states
  - -L is a set of labels (e.g. actions, or conditions)
  - $\rightarrow \subseteq S \times L \times S$  is a transition relation
- $\bullet \ \ {\rm We write} \ p \xrightarrow{a} q \ {\rm to} \ {\rm mean} \ (p,a,q) \in \rightarrow$
- An LTS is not a DFA or NFA



 $\bullet$  A simulation is a binary relation  $R \subseteq S \times S$  such that if  $(p,q) \in R$  then whenever

$$p \xrightarrow{a} p'$$

then

$$q \xrightarrow{a} q'$$

and

 $(p',q')\in R$ 

- p and q are similar, written  $p \preceq q$  if there is a simulation relation R such that  $(p,q) \in R$ 

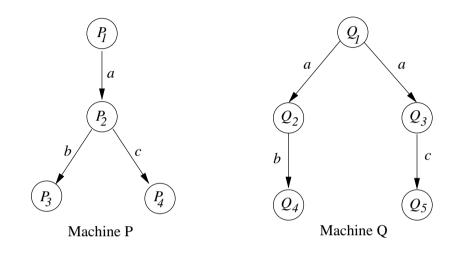


• A bisimulation is a binary relation  $R \subseteq S \times S$  such that if  $(p,q) \in R$  then

— whenever 
$$p\xrightarrow{a}p'$$
 then  $q\xrightarrow{a}q'$  and  $(p',q')\in R$  , and

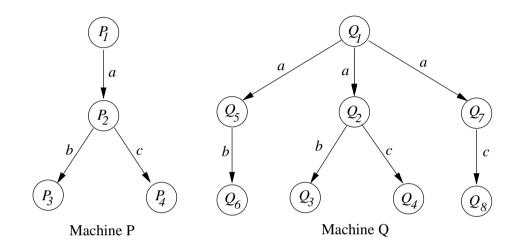
- whenever  $q \xrightarrow{a} q'$  then  $p \xrightarrow{a} p'$  and  $(p',q') \in R$
- p and q are bisimilar, written  $p \sim q$  if there is a bisimulation relation R such that  $(p,q) \in R$
- $\bullet~$  Note: If  $p \preceq q$  and  $q \preceq p$  then it is not necessarily the case that  $p \sim q$





- $Q_1 \preceq P_1$  because there is a simulation R s.t.  $(Q_1, P_1) \in R$
- $R = \{(Q_1, P_1), (Q_2, P_2), (Q_3, P_2), (Q_4, P_3), (Q_5, P_4)\}$





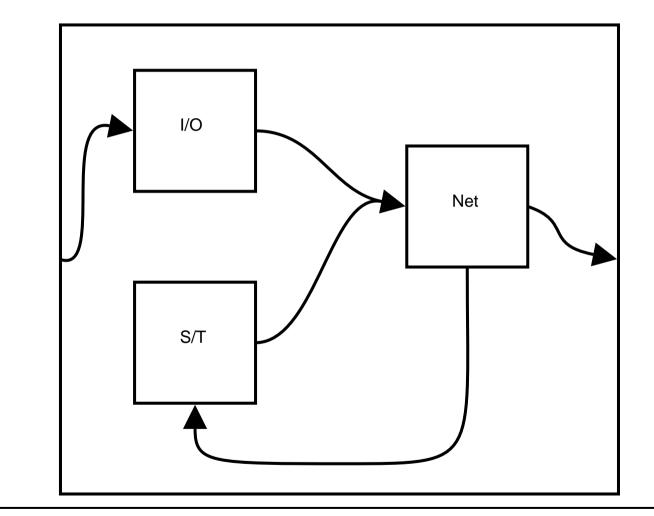
•  $Q_1 \preceq P_1$  and  $P_1 \preceq Q_1$  but  $Q_1 \not\sim P_1$ 



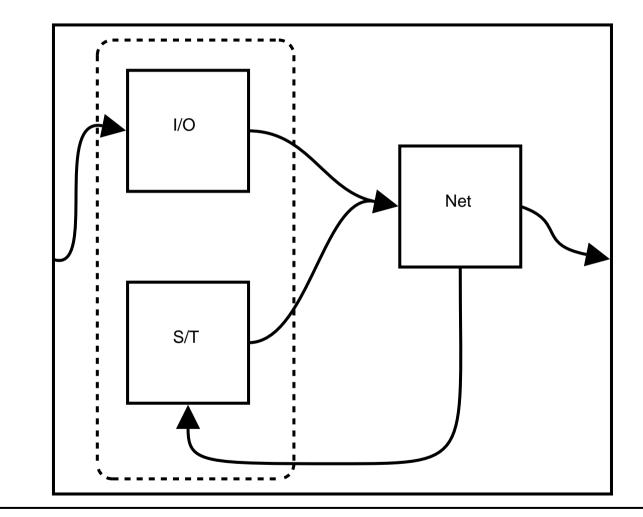


- At the Network level
  - Matching interfaces
  - Graph-homomorphism



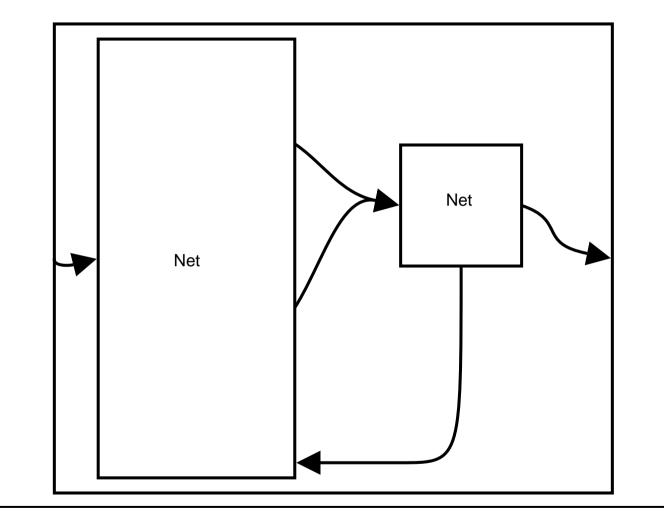








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### Progress

✓ Motivation

- ✔ Hierarchies of System Specification
- ✓ System morphisms
  - $\checkmark$  System morphisms at the same level of specification
  - $\pmb{\times}$  System morphisms at different levels of specification
- **★** Analysis and verification
- **✗** Category Theory: the meta-theory



- Structural
- Behavioural
- Mixed



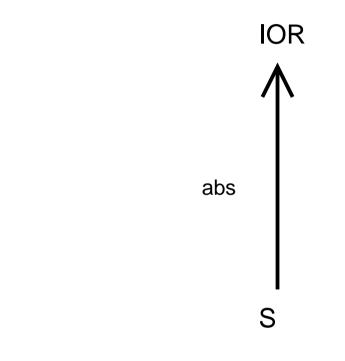
- Structural
  - Between Observation Frame and Network
  - Between Networks
- Usually given by an homomorphism
- Answers the question:

"Can I plug-in this component in that network?"



- Behavioural
  - Between I/O relation and State/Transition
- Answers the questions:
  - "What is the behaviour of this system?"
  - $\mathsf{and}$
  - "Given this I/O relation, does that system satisfy it?"

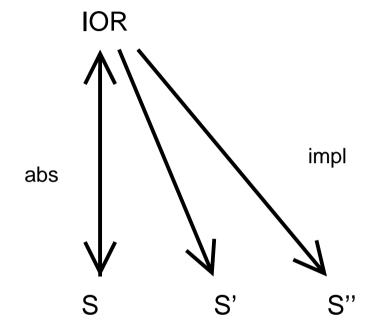




• S satisfies IOR



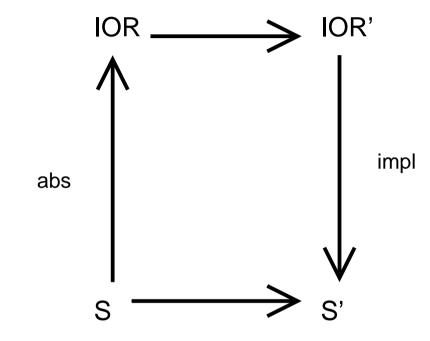




• IOR is implemented by S, S', S'', etc.



System morphisms at different levels of specification



- $IOR(S) \subseteq IOR(S')$  if and only if  $S \preceq S'$
- IOR(S) = IOR(S') if and only if  $S \sim S'$

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## System morphisms at different levels of specification

- Mixed
  - Between I/O relation and Observation Frame
  - Between State/Transition and Observation Frame
  - Between State/Transition and Network



# System morphisms at different levels of specification

- Bisimilarity: observational/behavioural equivalence
- $\bullet$  If two systems P and Q are equivalent then no observer should be able to distinguish between them
- If a part P of a composite system (network) C[P] is replaced by another equivalent part Q, then the resulting system C[Q] should behave in the same way
- For all contexts C[-], if  $P\sim Q$  then  $C[P]\sim C[Q]$
- $\bullet \sim$  should be a congruence



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# Analysis and verification

- Analysis: reasoning about systems and formalisms
- Properties:
  - System specific
  - Formalism specific
  - General



# Analysis and verification

- Techniques for establishing system specific properties:
  - Manual: by inspection, by formal analysis
  - Automatic: model-checking
- Modal logics: expressing system properties
  - Temporal logics: LTL, CTL, CTL\*, etc.
  - Epistemic logics
  - etc.
- Given a system (and a state,) and some formula, determine if it is satisfied or not



# Analysis and verification

- Techniques for establishing formalism specific and general properties
  - Manual: by general induction (on the system structure, on the proof of the property, etc.)
  - Automatic: theorem-proving



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- A framework for expressing and relating different mathematical concepts.
- A *category* is a mathematical structure that represents a family of *objects* (mathematical structures) **and** their relationships (morphisms.)

- Examples:
  - **Set**: the category of sets and functions
  - **Rel**: the category of sets and relations
  - Mon: the category of monoids and monoid homomorphisms
  - **Pre**: the category of preorders and monotonic functions
  - $\ensuremath{\text{Vec}}$ : vector spaces and linear transformations
  - Top: topological spaces and continuous functions
  - Graph: the category of graphs and graph-homomorphisms
  - **ST**: state/transition systems and simulations
  - **Prog**: data-types and programs



- A *functor* is a map between two categories.
- A *natural transformation* is a map between two functors (that go the same way.)
- An *adjunction* is a relation between two functors (that go in opposite ways.)



• Abstraction and refinement are adjoint functors



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# The end

# TaDa!

